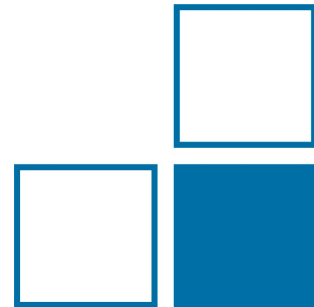


# Blockchain and eVoting

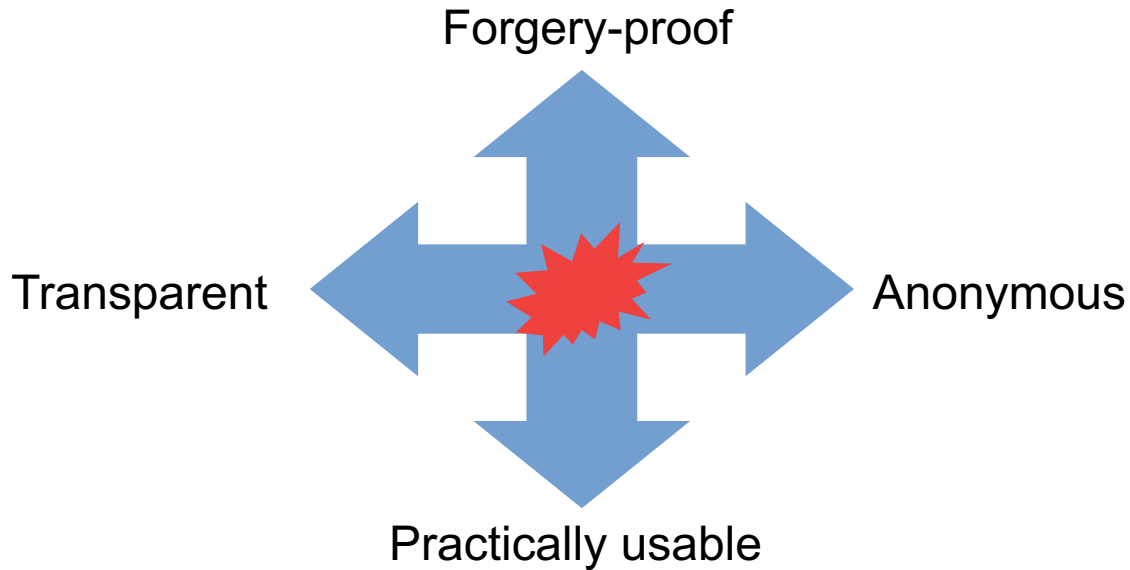
Daniel Peters, PTB

WG 8.52 Embedded Metrological Systems



Electronic voting is an **online** process in which **registered** voters cast their vote from an electronic device and transmit it via the Internet to an electronic ballot box (or **bulletin board**).

- **Construct a network in which voters can anonymously change their ballot even after they already cast it, before the deadline passed.**



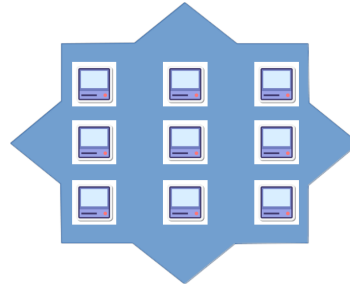
IT mechanisms for secure e-Voting:

- **(Blind) Signatures**
- **Distributed Ledger Technology**
- **Mix Networks (Onion-encryption)**
- **Zero-knowledge Proofs**
- **Identity Based Cryptography**
- **(a)symmetric Encryption**
- **Public Key Infrastructure**
- **Hash functions**
- **Homomorphic Encryption Schemes**
- **Multi-party computations**

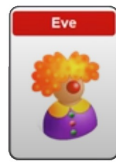
Central Authority

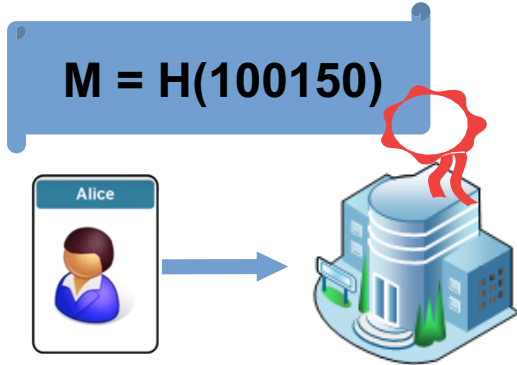


Anonymization Network



Public Database





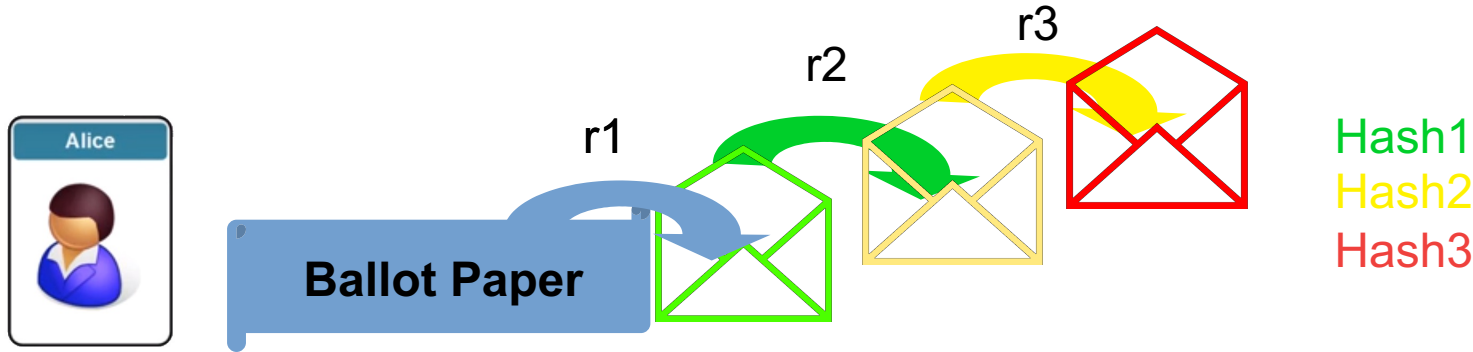
Alice wants her  $m$  to get signed by CA without disclosing it.

- 1) Alice uses the public key  $A$  from CA to transform the message  $m$  to an unreadable form by using a random number  $k$ .
- 2) As a result of mathematically combining  $m$ , public key  $A$  and  $k$  she generates an encrypted message  $m^*$  and sends  $m^*$  to CA
- 3) CA deletes Alice out of the list of allowed voters to avoid double votings
- 4) CA uses  $m^*$  and its private Key from  $A$  to create a signature on  $m^*$  obtaining  $s^*$  and sends the message to Alice securely
- 5) Alice uses  $k$  to revert her encryption on  $s^*$  and gets a signature  $s$  from CA for  $m$  without disclosing  $m$  to CA.

We make use of two different tokens, called:

- 1) The Network Usage Token (NUT)
- 2) The Initial Voting Token (IVT)

- These are all blind signed hashes by the CA
- Voter generates  $i + b + n$  different random numbers and their hashes.
- CA generates a new key pair ( $pk\_NUT$  ,  $sk\_NUT$  ) for every time slot

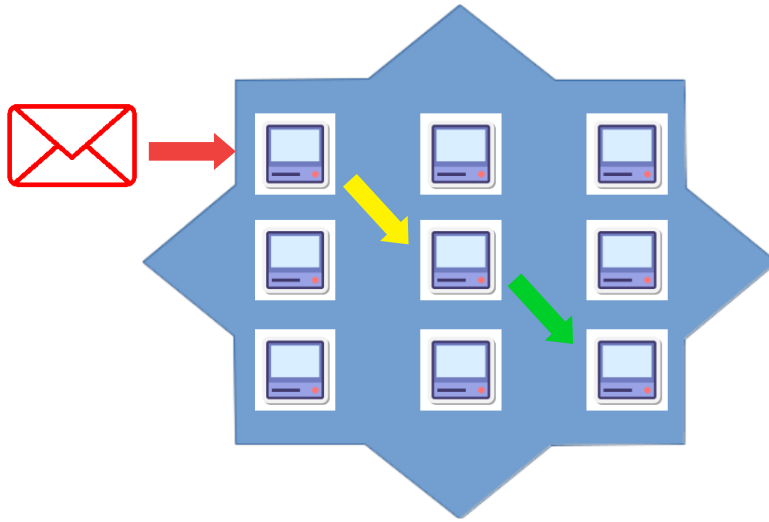


1. Alice uses the public key of node1 to encrypt her Ballot and a random number ( $r_1$ ) chosen by Alice (green envelope)
2. She creates a hash value of „green envelope“ and saves it together with  $r_1$
3. Alice uses public key of node 2 to package the „green envelope“ in yellow one adding another random number and noticing the hash of „yellow envelope“
4. She does the same for the red envelope using node 3 public key



**To open an envelope the corresponding node private key is needed, which is only known to the node itself**

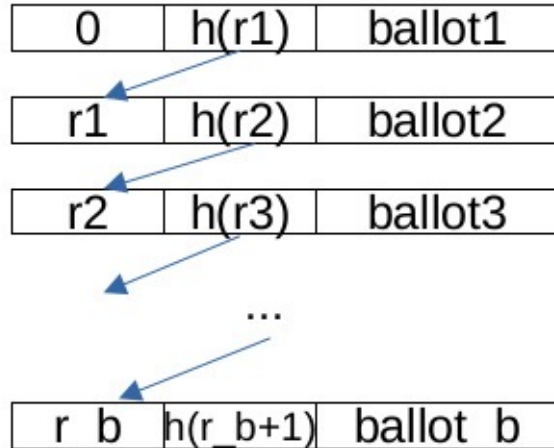




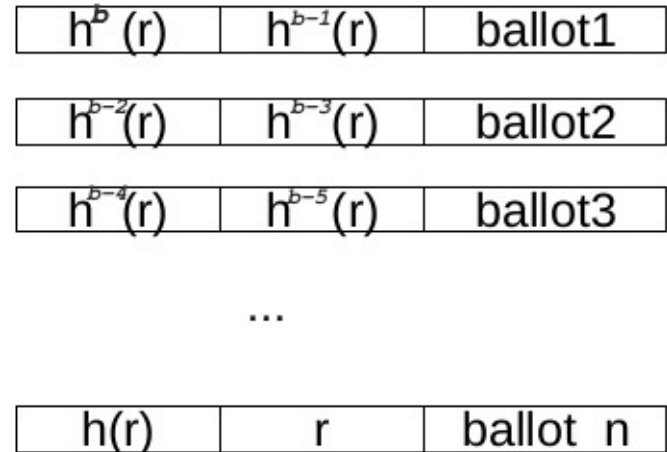
**Every step is  
synchronized by the CA  
by defining time slots**

- 1) MixServer1 notices the „yellow envelope“ on the public database
- 2) MixServer1 extracts the „yellow envelope“ and writes the „green envelope“ to the database
- 3) MixServer 1 generates a zkProof to show that mixing is correct
- 4) MixServer2 does the same ...

IVT signed

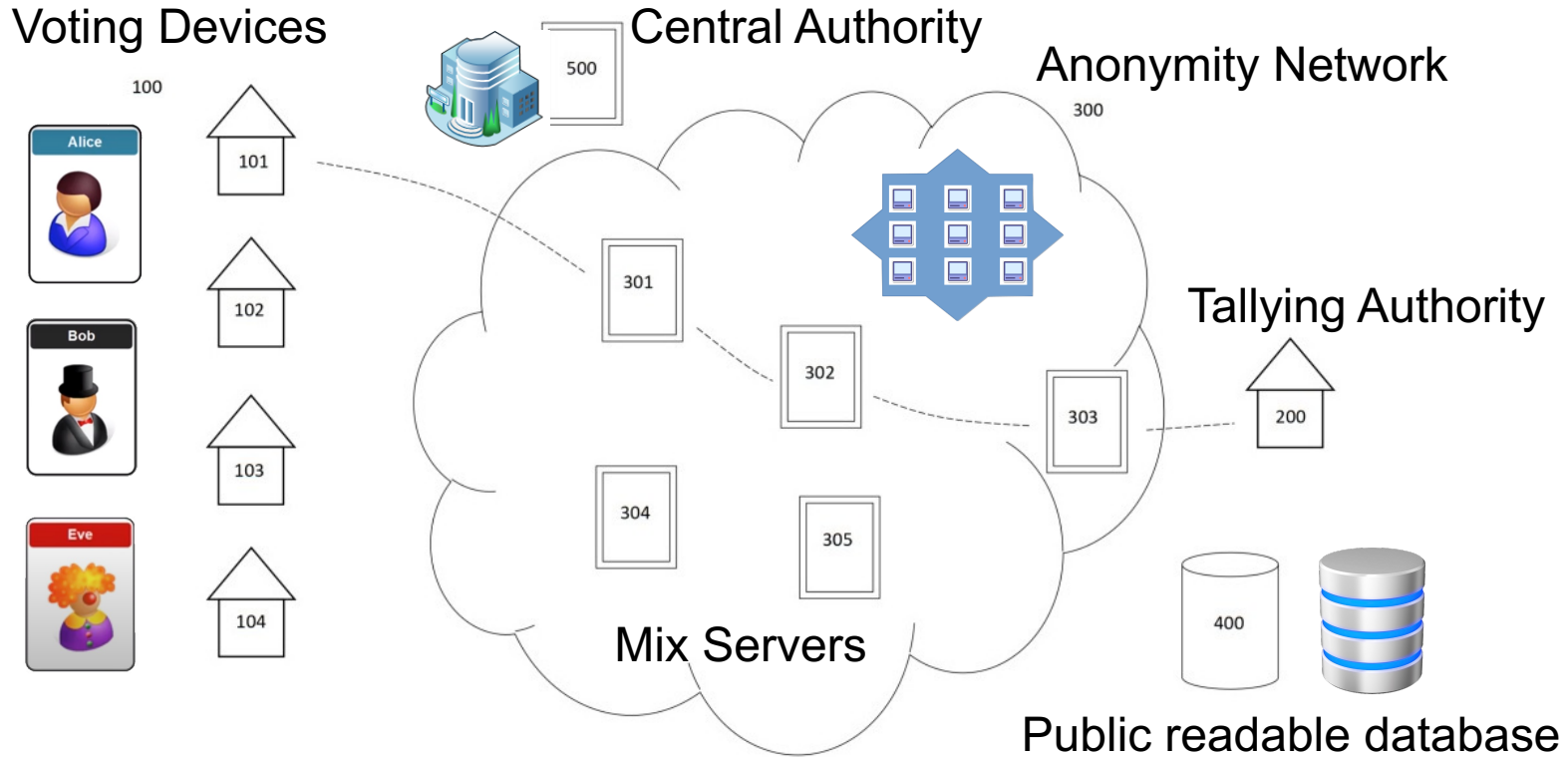


IVT signed



Ballots before union  
encryption

# PTB Our network topology

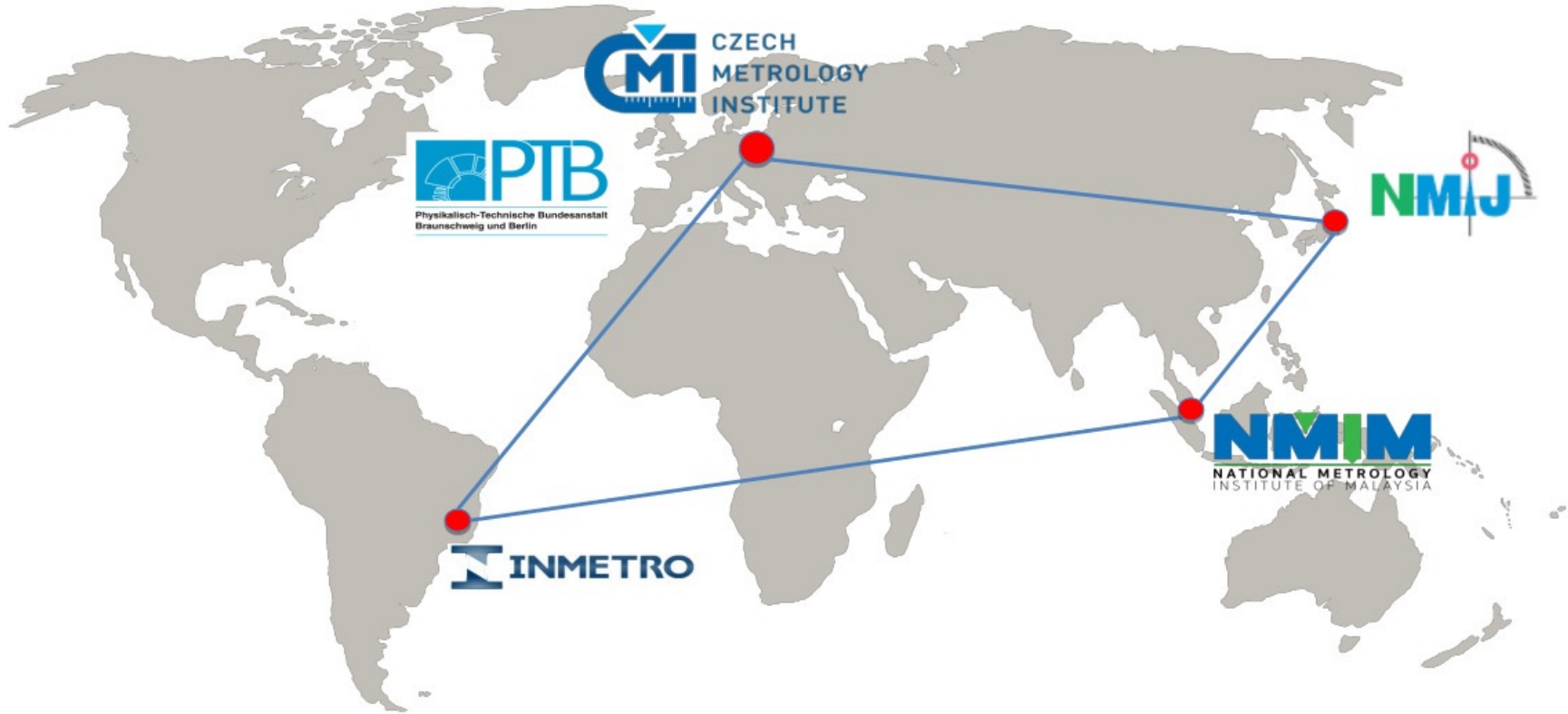


- Is accessible all the time
- Checks if voters are permitted to vote
- Generates blind signatures for IVT and NUT
- Synchronizes the anonymity network
- Verifies the zkProofs

- The MixNet is synchronized by the CA and every stepped is logged
- The synchronization takes place by using time-frames
- Logged in a public readable database which is permissioned
- Only the CA has write-access to it

- An easy to use graphical user interface
- The blind signature implementation, for the communication with the CA
- Secure random number creation
- ID checking on the device
- Communication with the anonymity network
- Choosing the order of the MixServers, through which the ballot should be onion-encrypted and sent

- 1) The tallying authority starts a new election through the CA
- 2) A voter identifies herself to the CA, and the CA gives authorization through blind signed tokens
- 3) The voter uses the anonymity network to submit her onion-encrypted vote. She authorizes herself for using the anonymity network through the signed NUT
- 4) The nodes in the anonymity network send all their batches of votes that they received in a time slot mixed with a zkProof to the CA. The CA checks the NUT tokens (from the entry MixServers) and the zkProofs for every step.
- 5) The CA puts the next layer of encryption onto the public database for the next MixServer
- 6) This is repeated until only the completely decrypted ballot paper is left
- 7) Because of the hash properties, voters can change their mind by sending new ballots that are linked with the old ones before the deadline has passed.







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